

Representability and the YONEDA-Lemma.

Def. a) Let $F: C \rightarrow \text{Set}$ be a (co- or contravariant) functor with C a locally small category.

A **representation** for F is a selection of an object $c \in C$ with a specified natural isomorphism

$$C(c, -) \cong F, \text{ if } F \text{ is covariant, or} \\ C(-, c) \cong F, \text{ if } F \text{ is contravariant.}$$

If such a c exist, then we say that F is **representable** and **represented by c** .

- b) Let C be a category and $c \in C$ an object. We say that c is
- **initial** if the constant functor $*$: $C \rightarrow \text{Set}$ that sends every object to the singleton set is represented by c .
 - **terminal** if the constant functor $*$: $C^{\text{op}} \rightarrow \text{Set}$ is represented by c .

remark. • c is initial $\Leftrightarrow C(c, a)$ is a singleton set for all $a \in C$.
• c is final $\Leftrightarrow C(a, c) \dashv \vdash$.

Examples. a) $1_{\text{Set}}: \text{Set} \rightarrow \text{Set}$ (identity functor) is represented by the singleton set 1 : $\text{Set}(1, X) = X^1 \cong X$ for every set X

Let $f: X \rightarrow Y$ be a morphism (= map) of sets:

$$\begin{array}{ccc} \text{Set}(1, X) \cong X & & (1 \mapsto x) \mapsto x \\ f_* \downarrow & \downarrow f & \downarrow \\ \text{Set}(1, Y) \cong Y & & (1 \mapsto f(x)) \mapsto f(x) \end{array}$$

b) $U: \text{Grp} \rightarrow \text{Set}$ (forgetful functor) is represented by $\mathbb{Z} \in \text{Grp}$.

$\varphi: G \rightarrow H$ group homomorphism:

$$\begin{array}{ccc} \text{Grp}(\mathbb{Z}, G) \cong UG & & (1 \mapsto g) \mapsto g \\ \varphi_* \downarrow & \downarrow \varphi & \downarrow \\ \text{Grp}(\mathbb{Z}, H) \cong UH & & (1 \mapsto \varphi(g)) \mapsto \varphi(g) \end{array}$$

(this is an isomorphism because every homomorphism $\mathbb{Z} \rightarrow G$ is determined by the image of $1 \in \mathbb{Z}$.)

c) R unital ring. $U: \text{Mod}_R \rightarrow \text{Set}$ is represented by $R \in \text{Mod}_R$.

$\varphi: M \rightarrow N$ R -module homomorphism:

$$\begin{array}{ccc} \text{Mod}_R(R, M) \xrightarrow{\cong} UM & (1 \mapsto m) \mapsto m & \\ \varphi_* \downarrow & \downarrow \varphi & \downarrow \\ \text{Mod}_R(R, N) \xrightarrow{\cong} UN & (1 \mapsto \varphi(m)) \mapsto \varphi(m) & \end{array}$$

d) $\text{Path}: \text{Top} \rightarrow \text{Set}$ is represented by $I = [0, 1] \subseteq \mathbb{R}$,

$\text{Loop}: \text{Top}_* \rightarrow \text{Set}$ ——— S^1 (with fixed base point)

by definition: $\text{Top}(I, X) \cong \text{Path } X \quad \forall X \in \text{Top},$
 $\text{Top}_*(S^1, Y) \cong \text{Loop } Y \quad \forall Y \in \text{Top}_*.$

e) (CURRYing; contravariant example)

$\text{Set}(- \times A, B): \text{Set}^{\text{op}} \rightarrow \text{Set}$ is represented by $B^A = \text{Set}(A, B)$

$\text{Set}(X \times A, B) \cong \text{Set}(X, B^A), \quad \forall X \in \text{Set}$

$(\varphi: X \times A \rightarrow B) \mapsto (\varphi^c: X \rightarrow B^A, x \mapsto (\varphi_x^c: A \rightarrow B, a \mapsto \varphi(x, a)))$

(Inverse: $(\psi: X \rightarrow B^A) \mapsto (\psi_c: X \times A \rightarrow B, (x, a) \mapsto [\psi(x)](a))$

f) Let $N \trianglelefteq G$ be a normal subgroup of the group G and $\pi: G \rightarrow G/N$ the canonical epimorphism.

$F: \text{Grp} \rightarrow \text{Set}$

$H \mapsto \{\varphi \in \text{Grp}(G, H) \mid N \leq \ker \varphi\} \in \text{Grp}(G, H)$

$(\psi: H \rightarrow H') \mapsto (\psi_*: FH \rightarrow FH', \varphi \mapsto \psi \circ \varphi)$

where $\ker \psi \circ \varphi \geq \ker \varphi \geq N$ because for $g \in \ker \varphi$

$$\psi \circ \varphi(g) = \psi(e_H) = e_{H'}.$$

Universal property of quotient group:

F is represented by $G/N \in \text{Grp}$.

If $\varphi: H \rightarrow H'$ is a group homomorphism, we obtain

$$\begin{array}{ccc} \text{Grp}(G/N, H) \xrightarrow{\cong} FH & \tilde{\varphi}: G/N \rightarrow H & \varphi: G \rightarrow H \\ \varphi_* \downarrow & \downarrow \varphi_* & \downarrow \varphi_* \\ \text{Grp}(G/N, H') \xrightarrow{\cong} FH' & \varphi_* \tilde{\varphi}: G/N \rightarrow H' & \varphi_* \varphi: G \rightarrow H' \\ [g] \mapsto \varphi_* \varphi(g) & [g] \mapsto \varphi_* \varphi(g) & g \mapsto \varphi_* \varphi(g) \end{array}$$

given by $(\tilde{\varphi}: G/N \rightarrow H) \mapsto (\varphi = \tilde{\varphi} \circ \pi: G \rightarrow H).$

($\ker \tilde{\varphi} \geq \ker \pi = N$ by definition)

The inverse is $\psi \in FH \mapsto (\tilde{\psi}: G/N \rightarrow H, [g] \mapsto \psi(g)),$
 which is well defined as $[g] = [h] \Leftrightarrow gh^{-1} \in N \leq \ker \psi,$
 so $\psi(g)\psi(h)^{-1} = 1 \Rightarrow \psi(g) = \psi(h).$

"old version"

$\varphi: G \rightarrow H$ Grp hom.

$N \leq \ker \varphi$

$\Rightarrow \exists! \tilde{\varphi}: G/N \rightarrow H,$

$\varphi = \tilde{\varphi} \circ \pi$

$$\begin{array}{ccc} G & \xrightarrow{\varphi} & H \\ \pi \downarrow & \nearrow \exists! \tilde{\varphi} & \\ G/N & & \end{array}$$

Problem. Given a functor $F: C \rightarrow \text{Set}$, what information is necessary to define a natural isomorphism $\alpha: C(c, -) \cong F$ - or more generally a natural transformation $\alpha: C(c, -) \Rightarrow F$?

Example. ordinal category ω ($0 \rightarrow 1 \rightarrow 2 \rightarrow \dots \rightarrow n \rightarrow \dots$).
 $F: \omega \rightarrow \text{Set}$ is a family $(F_n)_{n \in \omega}$ of sets with functions $f_{n, n+1}: F_n \rightarrow F_{n+1}$, $n \in \omega$. $\omega(k, -): \omega \rightarrow \text{Set}$, represented by $k \in \omega$, maps $n \in \omega$ to the empty set, if $n \leq k$, or to a singleton set $*$, else.

If $\alpha: \omega(k, -) \Rightarrow F$ is a natural transformation, we obtain:

$$\begin{array}{ccccccc} \emptyset & \rightarrow & \emptyset & \rightarrow & \dots & \rightarrow & \emptyset & \rightarrow & * & \rightarrow & * & \rightarrow & \dots \\ \alpha_0 \downarrow & & \alpha_1 \downarrow & & & & \alpha_{k-1} \downarrow & & \alpha_k \downarrow & & \alpha_{k+1} \downarrow & & \\ F_0 & \xrightarrow{f_{0,1}} & F_1 & \rightarrow & \dots & \rightarrow & F_{k-1} & \xrightarrow{f_{k-1,k}} & F_k & \xrightarrow{f_{k,k+1}} & F_{k+1} & \rightarrow & \dots \end{array}$$

the α_n , $n \leq k$, carry no information and the α_n , $n \geq k$, are equivalent to a single element from each F_n which we will also call α_n . Then naturality implies

$$\alpha_{n+1} = f_{n,n+1}(\alpha_n) = f_{n,n+1}(f_{n-1,n}(\dots(f_{k+1,k}(\alpha_k))\dots)) \quad \forall n \geq k.$$

Thus α is completely determined by $\alpha_k = \alpha_k(1_k) \in F_k$ with $1_k \in \omega(k, k)$.

Theorem. (YONEDA lemma)

Let C be a locally small category and $F: C \rightarrow \text{Set}$ a functor. Then, for every $c \in C$ there is a bijection

$$\text{Hom}(C(c, -), F) \cong F_c,$$

that associates $\alpha: C(c, -) \Rightarrow F$ to $\alpha_c(1_c) \in F_c$.

Furthermore, this correspondence is natural in c and F .

This also implies that $\text{Hom}(C(c, -), F)$ is a set.

Proof. Obviously,

$$\Phi: \text{Hom}(C(c, -), F) \rightarrow F_c, \alpha \mapsto \Phi(\alpha) = \alpha_c(1_c)$$

is a function. We now need to find an inverse for it:

Let $x \in F_c$ be arbitrary. We now desire a natural transformation $\mathbb{I}(x): C(c, -) \Rightarrow F$, i.e. for every $f \in C(c, d)$ we want, that

$$\begin{array}{ccc} \text{the square} & C(c, c) & \xrightarrow{\Phi(x)_c} F_c \\ & \downarrow f_* & \downarrow Ff \\ & C(c, d) & \xrightarrow{\Phi(x)_d} F_d \end{array}$$

commutes. For $1_c \in C(c, c)$, this implies

$$\hookrightarrow \Phi(x)_d(f) = Ff(\Phi(x)_c(1_c)) \quad \downarrow$$

As Ψ is supposed to be an inverse of Φ , we conclude

$$\Phi(\Psi(x)) = \Psi(x)_c(1_c) \stackrel{!}{=} x \in F_c$$

Thus we define

$$\Psi: F_c \rightarrow \text{Hom}(C(c, -), F), \quad \Psi(x)_d(f) := Ff(x).$$

Now let $g: d \rightarrow c$ be an arbitrary morphism in C . We want the following square to commute:

$$\begin{array}{ccc} C(c, d) & \xrightarrow{\Phi(x)_d} & F_d \\ g_* \downarrow & & \downarrow Fg \\ C(c, c) & \xrightarrow{\Phi(x)_c} & F_c \end{array}$$

For $f \in C(c, d)$ arbitrary, we obtain:

$$\begin{aligned} \Phi(x)_c \circ g_*(f) &= \Phi(x)_c(g \circ f) = F(g \circ f)(x) \stackrel{F \text{ functor}}{=} Fg \circ Ff(x) \\ &= Fg(\Phi(x)_d(f)) = Fg \circ \Psi(x)_d(f) \end{aligned}$$

and as such $\Psi(x)$ is natural and Ψ is well-defined.

Lastly we want to show that Ψ is a left inverse of Φ (right invertibility: $\Phi(\Psi(x)) = \Psi(x)_c(1_c) = x$ by definition).

Let $\alpha: C(c, -) \Rightarrow F$ be a natural transformation. Then:

$$\Psi(\Phi(\alpha)) = \Psi(\alpha_c(1_c)).$$

By the commutativity of the square

$$\begin{array}{ccc} C(c, c) & \xrightarrow{\alpha_c} & F_c \\ f_* \downarrow & & \downarrow Ff \\ C(c, d) & \xrightarrow{\alpha_d} & F_d \end{array} \quad (*)$$

we obtain

$$\Psi(\Phi(\alpha))_d(f) = Ff(\alpha_c(1_c)) = \alpha_d(f)$$

and thus $\Psi(\Phi(\alpha)) = \alpha$.

Naturality in F :

Given a natural transformation $\beta: F \Rightarrow G$, we want to show that the corresponding element to $\beta \circ \alpha: C(c, -) \Rightarrow F \Rightarrow G$ in G_c is the image of that corresponding to $\alpha: C(c, -) \Rightarrow F$

in F_c under $\beta_c: F_c \rightarrow G_c$:

$$\begin{array}{ccc} \text{Hom}(C(c, -), F) & \xrightarrow[\cong]{\Phi_F} & F_c \\ \beta_* \downarrow & & \downarrow \beta_c \\ \text{Hom}(C(c, -), G) & \xrightarrow[\cong]{\Phi_G} & G_c \end{array}$$

This diagram obviously commutes as:

$$\hookrightarrow \Phi_G(\beta_* \alpha) = (\beta_* \alpha)_c(1_c) = \beta_c(\alpha(1_c)) = \beta_c(\Phi_F(\alpha)) \quad \checkmark$$

Similarly, given a morphism $f: c \rightarrow d$ in C , we want to show that the corresponding element to $\alpha \cdot f^*: C(d, -) \Rightarrow C(c, -) \Rightarrow F$ in F_d is the image of the corresponding element to $\alpha: C(c, -) \Rightarrow F$ in F_c under $Ff: F_c \rightarrow F_d$:

$$\begin{array}{ccc} \text{Hom}(C(c, -), F) & \xrightarrow[\cong]{\Phi_c} & F_c \\ (f^*)^* \downarrow & & \downarrow Ff \\ \text{Hom}(C(d, -), F) & \xrightarrow[\cong]{\Phi_d} & F_d \end{array}$$

We can see:

$$\checkmark Ff(\Phi_c(\alpha)) = Ff(\alpha_c(1_c)) = (\alpha \circ f^*)_d(1_c) = \Phi_d((f^*)^* \alpha) \quad \hookrightarrow$$

as

$$\begin{array}{ccccc} C(d, d) & \xrightarrow{f^*} & C(c, d) & \xrightarrow{\alpha_d} & F_d \\ 1_d & \mapsto & f & \mapsto & \alpha_d(f) \end{array}$$

and (*) implies

$$(\alpha \circ f^*)_d(1_c) = \alpha_d(f) = Ff(\alpha_c(1_c))$$

as desired. $\square$

Remark. $(c, F) \in C \times \text{Set}^C$, where Set^C is the category of functors $C \rightarrow \text{Set}$ with natural transformations between them as morphisms. $\text{ev}: C \times \text{Set}^C \rightarrow \text{Set}$, $(c, F) \mapsto F_c$ is a bifunctor

$$\begin{array}{ccc} C' & \xrightarrow{y} & \text{Set}^C \\ \downarrow f & & \uparrow f^* \\ C & \xrightarrow{y} & \text{Set}^C \\ d & \mapsto & C(d, -) \end{array}$$

possible size issue here

$$\begin{array}{ccc} \text{Hom}(y(-, -)) : C \times \text{Set}^C & \xrightarrow{y \times 1_{\text{ev}}} & (\text{Set}^C)^{op} \times \text{Set} \xrightarrow{\text{Hom}} \text{Set} \\ (c, F) & \mapsto & (C(c, -), F) \mapsto \text{Hom}(C(c, -), F) \end{array}$$

• YONEDA lemma is equivalent to

$$\begin{array}{ccc} C \times \text{Set}^C & \xrightarrow{\text{Hom}(y(-, -))} & \text{Set} \\ \downarrow \Phi & & \downarrow \text{ev} \\ C \times \text{Set}^C & \xrightarrow[\cong]{} & \text{Set} \end{array}$$